

UNIVERSITY OF ILLINOIS  
DIGITAL COMPUTER

LIBRARY ROUTINE MA1 - 104

**TITLE** Matrix Multiplication Floating Decimal Auxiliary

**TYPE** Floating Decimal Auxiliary; uses interpretive orders and control registers 0, 1, and 2.

**NUMBER OF WORDS** 21 without punchout -- 26 with punchout

**DURATION** Without punchout,  $9 + 3n(2p + 2m + 11) + 10mp$  millisec. plus input of elements.  
With punchout,  $9 + 3n(2p + 2m + 11) + 19mp$  milliseconds plus input and output time.

**ACCURACY** Usually 9 decimals, but depends on condition of matrices i.e. inversely proportional to the range of the elements in any one matrix.

**TEMPORARY STORAGE** 0, 1, 2, 3 mp words specified by SN, n words specified by SS, np words specified by SK

**PARAMETERS** S3 - SJ

3 = (Address of Floating Accumulator)  $\times 2^{-39}$   
4 = (address of Code A1)  $\times 2^{-39}$   
5 =  $np \times 2^{-39}$   
6 =  $n \times 2^{-39}$   
7 =  $p \times 2^{-39}$   
8 =  $mp \times 2^{-39}$   
9 =  $e \times 2^{-39}$   
K =  $a_{00} \times 2^{-39}$   
S =  $b_{10} \times 2^{-39}$   
N =  $c_{00} \times 2^{-39}$

Where: B is an (m by n) matrix, A is an (n by p) matrix, and C = BA is an (m by p) matrix.  $a_{00}$  is the location of the first element of A.  $b_{10}$  is the location of the first element of any row i of B.  $c_{00}$  is the location of the first element of C = BA. e is the location of the left hand order to which control is transferred upon completion of this routine. If punch out is desired, set e = 21L in this routine.

**MAXIMUM MATRICES** Square matrices maximum  $m = n = p = 20$ . Always,  $mp + n + np \leq 825$

**DESCRIPTION** This auxiliary routine performs matrix multiplication of two matrices B = m by n matrix and A = n by p matrix to obtain BA = C = m by p matrix in floating decimal form. The elements of C are stored in row form into successive locations beginning with SN. The routine is entered by an interpretive 8J order or by an immediate 26 or 24 xN directive.

**METHOD**

The method used is the same as that of the regular matrix multiplication routine, Code M1 except that no terminating symbol proceeds each row.

**PUNCHING AND SCALING**

To insure that no element  $c_{ik}$  of C exceeds  $10^{64}$ ,  $|a_{jk}|$  and  $|b_{ij}|$  should be  $\leq \sqrt{10/n} \times 10^{32}$ . The elements of A and B are punched in row form for regular floating decimal input. The elements of A are punched first and then the elements of B. When started, the routine inputs all of A and then inputs B one row at a time.

**NOTE**

- (1) If the elements in any one matrix vary over an excessive exponential range ( $> 10$ ) the subproducts ( $b_{io} a_{jk}$ ) may induce an error. e.g. If subproducts of an element  $c_{ik}$  are  $10^{-20}$ ,  $+ 10^{10}$ ,  $+ 10^{15}$ ,  $- 10^{10}$ ,  $- 10^{15}$ ,  $- 10^{20}$ , in that order,  $c_{ik}$  will be represented by  $-10^{-20}$  since adding first  $10^{-20} + 10^{10}$  gives  $10^{10}$  in the machine, and by adding  $10^{15} - 10^{10} - 10^{15}$  to this we get 0. Then adding  $0 - 10^{-20}$  gives  $10^{-20}$  due to the range of the exponents exceeding 10.
- (2) If an element of A or B is represented by  $1/10 \times 10^p$ , it must be remembered that  $1/10$  has only 28 significant binary digits. Hence, for example, if subproducts of an element  $c_{ik}$  of C are  $(1/10) (4/10) \times 10^p$ , we may not get zero since the two terms in this sum will not be represented in the same way in the machine. The size of the error induced in such a case is proportional to the size of p.

DATE <u>7/23/53</u> Rt: <u>6/4/58</u>
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LOCATION	ORDER	NOTES	PAGE 1	MAI
0	41 3F 50 L		Clear 3 for storage	
1	26 S4 OK S5		Enter interpretive code Set loop to input np elements of A.	
2	88 F OS (SK)	From 3 by 3	Input $a_{jk}$ and store	
3	03 2L 1K S8	Yes ↓	Is A input? Set loop for counting elements of C.	
4	OK S6 88 F	From 13 <sup>r</sup> From 5 <sup>r</sup>	Set loop to input elements of B Input row i of B and store.	
5	OS (SS) 02 4L	Yes ↓	Is row i of B input?	
6	2K S7 OK S6	From 12 <sup>r</sup>	Set loop for p $\sum$ s using row i of B. Set loop for n products to sum.	
7	05 (SS) 87 (SK)	From 9 <sup>r</sup> By 15	$b_{ij}$ ( $b_{ij}^a a_{jk}$ )	
8	84 3F 8S 3F		$\sum b_{ij}^a a_{jk} = N(3)$	
9	8J 14L 03 7L	Yes ↓	Escape to 14L Is $\sum b_{ij}^a a_{jk}$ complete for column k of A?	
10	85 3F 1S (SN)	↓	Store $c_{ik}$	
11	8J 16L 12 12L	Yes ↓	Escape to 16L Has C been computed?	
12	8J S9 22 6L	Yes ↓ From 11 <sup>r</sup>	Escape when finished Have p sums been formed using row i of B?	
13	8J 19L 83 4L	Yes ↓	Then escape to 19L	
14	L5 7L L4 18L	From 9	Increase address of $a_{jk}$ by p	
15	42 7L 26 29S4	From 18,20	Reenter interpretive code	
16	41 3F L5 7L	From 11	Clear 3 for storage	

LOCATION	ORDER		NOTES	PAGE 2 MA1
17	10 1L		Reset address of $a_{jk}$ for next summation	
	14 43S4			
18	26 15L		Address used	
	00 S7			
19	41 S3	From 13	Clear S3 for 83 order.	
	15 20L			
20	26 15L		Address used.	
	00 SK			
21	32 21L	From 12	For Print out	
	50 21L			
22	26 S4		Enter Code A1	
	0K S8		Set cycle for $m_p$ elements of C.	
23	05 (SN)	From 24	$c_{ik}$	
	89 9F	By 24	Print $c_{ik}$ to 9 decimals	
24	03 23L	Yes	Have $m_p$ elements of C been printed?	
	8J 25L		Escape to stop.	
25	0F F	From 24		
	0F F			