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PURPOSE

The Raytheon 703 Input /Output Software (IOS) provides the user with a powerful and flexible file oriented input/output system. Features of the system include device reassignment, simultaneous I/O capability, centralized I/O monitor, end action, and I/O macros. IOS provides a common I/O system for all users and alleviates the programmer's need to understand the machine language I/O operations of the 703 computer, while still allowing full use of the 703 capabilities.

IOS is able to perform I/O with any standard 703 peripheral device through the use of the 703 priority interrupt system, which allows rapid response to the necessary I/O interrupts. To further minimize time spent doing I/O the software is file oriented; i. e. all information necessary to perform the I/O operation is contained in a File Input/Output Table (FIOT, see 1.3) which is readily available to the device drivers. IOS operates entirely in the automatic interrupt mode which enables processing concurrent with I/O operations. This allows the user to make more efficient use of the computer than he might otherwise provide for himself.

Several I/O macros are available in IOS. These macros reduce the effort required to program efficient input/output. The user need only become acquainted with macros OPEN, DOIO, and STAT to communicate with any standard 703 device.

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1.1 IOS FUNCTIONS

Several operations are available in IOS. These are:

1. Read
2. Write
3. Rewind
4. Backspace
5. Write end of file
6. Search for end of file
7. Write skip
8. Status
9. Verify disk
10. Punch leader

Not all functions are available for every type of I/O device. For example, one cannot rewind the teletype. Refer to the driver descriptions for applicable functions.

IOS is an input/output control system which provides a user the facility to perform I/O operations which are effectively independent of the actual physical device being used. This is achieved by allowing different physical devices to be substituted for I/O operations without requiring re-assembly or recompilation of the program. For example, a program can be debugged using a small set of test data punched on cards and read from the card reader. After debug, during production runs, the input can be changed to magnetic tape to process the actual data files.

A user performs I/O by requesting data transfers on logical units, which can be related to any physical device. In a typical I/O request to IOS, the user specifies a Logical Unit Number (LUN), the address of the I/O data buffer and the number of words to be or written.

The logical unit is associated with a particular driver which performs a data transfer on a specific I/O device. The user can, at load time, cause this association to change so that a different driver and, therefore, physical device is referenced by the I/O request. This driver/device association is kept in a Peripheral Equipment Assignment Table (PEAT, see 1.2).

The device number or LUN assignment is made at system creation time. The IOS user uses the established LUNs and is free to reassign a given LUN to one of a set of devices, through X-RAY control directives. (These directives are described in the EXEC documentation).

All information pertaining to an I/O operation is kept in an 8-word File Input/Output Table, (FIOT - see 1.3). This table details the specific operation for the I/O request, as well as the LUN, I/O buffer, and the number of words. FIOT also provides temporary storage area as needed by the device driver.

1.1.1 IOS Calls

Macros are communication links between the programmer and IOS. All calls go through IOS subroutines to the appropriate device drivers. They translate information from an argument calling sequence to a FIOT, perform the I/O operation, provide error checking and end-action.

1.1.2 IOS Drivers

An I/O driver is a subroutine designed to perform data transfers between a specific peripheral device and the computer. An I/O driver is never directly used by a programmer, but only indirectly via IOS macros. Thus, for example, a programmer requests data to be read from some logical unit. IOS macros pass the FIOT address to the drivers, and the driver then extracts necessary information to start the I/O operation. The driver then exits back to the user program.

1.1.3 IOS Interrupt Service

Once initiated, I/O operations are processed independent of the user's program. Data transfers for devices on the Direct Memory Access (DMA) channel (such as disk and magnetic tape) are performed automatically by the hardware. The interrupt service routine portion of each driver performs all data transfer operations for devices on the Direct Input/Output (DIO) channel (such as; paper tape reader, teletype, card reader, etc.) When the I/O operation is finished, then interrupt service routine for DIO and DMA devices perform all operations required to conclude the I/O operation. Device status and ending buffer address are stored in the FIOT. At this time an end-action subroutine may be executed, which permits any arbitrary operation to be performed as a part of the interrupt service function.

1.2 PERIPHERAL EQUIPMENT ASSIGNMENT TABLE (PEAT)

IOS processes input and output operations via logical equipment. Logical equipment is a functional classification of I/O devices; such as, symbolic input devices, listing devices, binary output devices, etc. Therefore, since program reference to I/O are by function, different devices may be assigned to a logical unit without requiring a compensating change in the program reference.

For example, the binary input device may be the high speed paper tape reader for one job and the teletype for the next, but neither the I/O monitor nor the users program need be reassembled. The device may simply be reassigned by an XRAY directive (refer to the X-RAY EXEC documentation).

The peripheral equipment assignment table contains the information about the correspondence between logical and physical units. IOS uses this table to select the appropriate drivers when input or output requests are made.

1.2.1 Logical Unit Definitions

The PEAT shown in Figure 1 contains more equipment than is included with a typical system, but it is needed to show how logical units function. The first 7 units are system logical units that are referred to by name in the remainder of this manual and in other manuals.

The definition of each of these names is listed below:

<u>LUN</u>	<u>Description</u>
0	SYSF <u>SY</u> Stem <u>F</u> ile. This is the unit from which the system library is read and written. In Figure 1 SYSF points to the disk.
1	SYSI <u>SY</u> Stem <u>I</u> ntput. SYSI is the unit for all directives input to system programs; such as the X-RAY EXEC or SYM II Assembler. In Figure 1 SYSI points to the teletype.

<u>LUN</u>	<u>Description</u>
2	PRIN <u>PR</u> imary <u>IN</u> put. PRIN is the unit for input to system programs where the input is neither a directive nor binary; i. e., SYM II accepts the input source text from PRIN. In Figure 1 PRIN points to the card reader.
3	LIST <u>LI</u> STing. System programs use LIST for tabular output other than error messages. SYM II outputs the assembly listing on the LIST device. In Figure 1 LIST points to the line printer.
4	BIN <u>BI</u> nary <u>IN</u> put. BIN is the system binary input unit. In Figure 1 BIN points to the card reader. Absolute or relocatable binary programs are input from this unit.
5	BOUT <u>BI</u> nary <u>OU</u> Tput. BOUT is the system binary output unit. In Figure 1 BOUT points to the card punch. SYM II outputs absolute or relocatable binary text to the BOUT unit.
6	SCR <u>SC</u> Ratch. System programs use SCR as temporary storage. SYM II uses SCR to store intermediate text. In Figure 1 SCR points to the disk.
7-11	LOGA, LOGB, LOGC, LOGD, LOGE The next 5 units, LOGA-E, are intended for users. Devices are divided into two classes; mass and non-mass. A mass device is one that can perform both read and write operations; i. e., disk, magnetic tapes, etc. Non-mass devices can perform only a read or write; i. e., card reader, card punch, etc. LOGA is the systems highest level non-mass input device. LOGB is the systems highest level non-mass output device. LOGC is the system mass unit. LOGD and E are assigned according to the assignment of other devices within the system. The user should look at the PEAT in his I/O monitor listing to determine how these units are assigned.
12	DUMMY This logical unit performs no I/O but allows programs to be checked out without a peripheral device being available.

The remainder of the logical units are assigned to the various physical devices contained in the system. The following are currently assigned:

<u>LUN</u>	<u>Assignment</u>
13	Console Teletype
14	Magnetic Tape Unit 0
15	Magnetic Tape Unit 1
16	Magnetic Tape Unit 2
17	Magnetic Tape Unit 3
18	Disk Unit 0
19	High Speed Paper Tape Punch
20	High Speed Paper Tape Reader
21	Card Reader
22	Card Punch
23	Line Printer
24	Multiplex Teletype
25	Disk Unit 1
26	Disk Unit 2
27	Disk Unit 3
28	Plotter

Physical assignments for all logical units are made when the monitor is assembled for each system. These assignments are determined by what units are physically present in the system. The user can determine what these initial (or standard) assignments are by referring to the PEAT table as it appears in the listing for his monitor. Following is an example of how the first 12 words of the PEAT might appear:

<u>LOC</u>	<u>Contents</u>	<u>Source Text</u>	<u>Comments Field</u>
3FD9	1200 3461 PEAT	BYTE ASYSF,0	0 SYSF
3FDA	0D00 3462	BYTE ASYSI,0	1 SYSI
3FDB	1500 3463	BYTE APRIN,0	2 PRIN
3FDC	1700 3464	BYTE ALIST,0	3 LIST
3FDD	1500 3465	BYTE ABIN,0	4 BIN
3FDE	1600 3466	BYTE ABOUT,0	5 OUT
3FDF	1200 3467	BYTE ASCR,0	6 SCR
3FE0	1500 3468	BYTE ALOGA,0	7 LOG A
3FE1	1600 3469	BYTE ALOGB,0	8 LOG B
3FE2	1200 3470	BYTE ALOGC,0	9 LOG C
3FE3	0E00 3471	BYTE ALOGD,0	10 LOG D
3FE4	0F00 3472	BYTE ALOGE,0	11 LOG E

To determine the assignment of a particular unit, find the unit name (or number) in the comment field. The left byte of the contents field contains the hex unit number of the assignment. Convert it to decimal and look further down the table for that number in the comments field. For example, the device assigned to BIN (logical unit 4) has unit number 15₁₆. (15₁₆ is the left byte of the pointer word corresponding to BIN). 15₁₆ converts to decimal 21 which is the unit number of the card reader.

1.2.2 PEAT Table Description

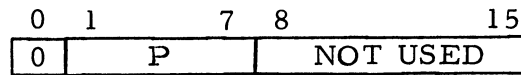
There are two types of entries in the PEAT. There are one word entries which define the correspondence between a logical unit and another logical unit, and there are two word entries which define the correspondence between a logical unit and a physical unit.

Entries within the PEAT are in logical unit number order; i. e., the first entry is for LUN 0, the second entry is for LUN 1, etc.

The table is in two sections. The first eleven entries, section 1, are always pointers to other LUNs. The remaining entries, section 2, are pointers to physical devices.

Numbers 0 - 11 may be considered logical unit numbers, whereas units 12 - 31 may be considered physical unit numbers, since they are not normally reassigned.

In section 1 there is one word per LUN, formatted as follows:

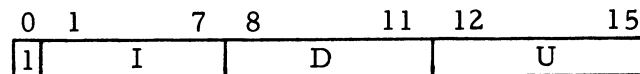


where:

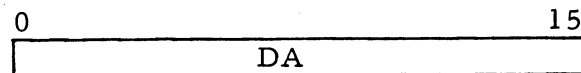
The 0 in bit 0 denotes a pointer to another unit.

P is a 7-bit field pointing to the location, relative to the top of the PEAT table, for the entry of the next LUN.

In section 2 there are two words per entry. If the unit has been reassigned, word one will have the same format as section 1, and word two will be unused. If the LUN has not been reassigned, the format is as follows:



Word 1



where:

The 1 denotes a pointer to a physical unit.

I is a 7-bit field used to denote the priority interrupt level for the device.

D is a 4-bit field used to denote the device code.

U is a 4-bit field used to denote the unit number of the device number; i. e., magnetic tape may be units 0, 1, 2, 3, etc.

DA is the device driver address for the LUN.

When IOS is seeking the physical unit corresponding to a logical unit, it will go from logical unit to logical unit until it finally arrives at an entry pointing to a physical unit (a 1 in bit position 0).

Daten

1.3 FILE INPUT/OUTPUT TABLE (FIOT)

The ^{Zweck} purpose of the FIOT is to provide user area for I/O operation ^{Beschreibung} description and temporary storage for device drivers. How each driver uses the temporary storage is described in the individual driver documentation. Each FIOT is 8 words in length and formatted as follows:

	0	1	7	11	15
FIOT 0	BB	IBA (<i>Büf Adr</i>)			
1	M	ANW (<i>Adr of Word count</i>)			
2				LUN	FC
3	S				
4	WA				
5	DS				
6	F	EAA			
7					

Word 0:

The Busy Bit (BB) is a one-bit field which is set to 1 while the I/O operation is being performed. It will be reset to 0 when the operation has completed. The Initial Buffer Address (IBA) is a 15-bit field which contains the word address of the first location of the data buffer to be transferred. *die die Wort-Adresse oder 1. Location des zu übertragenden Daten Puffers enthält*
 Word 1:

Adresse in der Information steht. Wort-Adresse

Mode (M) is a one-bit field which has different meanings for various physical devices and is documented with each device driver. (See 2.1). Anw is a 15-bit field which is the Address of a location which contains the Number of Words to be input or output.

Word 2:

Bits 0 through 6 are temporary storage. The Logical Unit Number (LUN) is a 5-bit field contained in bits 7 through 11. The Function Code (FC) is a 4-bit field contained in bits 12 through 15, and defines the type of operation that is to be performed. Certain function codes perform special functions on specific devices; these are described in the applicable driver documentation. Function codes 9 and B₁₆ will perform a read operation for every device capable of a read. Function code E₁₆ will perform a write operation on every device capable of a write. Codes 9 and B₁₆ are equivalent for every device ^{nicht für} except the teletype. To preserve device independence codes 9, B, and E should be used for all read and write operations.

Word 3:

Status is the status response word obtained from the device upon completion of each I/O operation. The response differs from device to device and is documented separately for each device. Refer to the appropriate driver documentation. For all devices bit 15 true means that an error has occurred during the I/O operation.

Word 4: nach Beendigung der I/O operation wird Wort 4 die Byte-Adresse des nächsten Bytes enthält.
 nach Beendigung der I/O operation wird Wort 4 die Byte-Adresse des nächsten Bytes enthält.

Working Address. Upon completion of the I/O operation word 4 will contain the byte address one greater than the buffer address of the last byte of data transferred. For certain devices this word contains the address of the next byte to be transferred throughout the I/O operation and may be monitored to determine the progress of the operation. Refer to individual driver documentation for details.

Word 5:

Word 5 serves as temporary storage for most drivers. For the disk driver, this word contains the Sector Address to be read or written. When performing device independent I/O, where the disk is a legitimate assignment for the unit, word 5 should be set to the desired sector. Then, if some other device is assigned, the disk sector specification will be ignored. The format for this word as a disk sector is described in the disk driver documentation.

Word 6:

Bit 0 of word six is used for record Format control. When bit 0 is false, those drivers which generate or read formatted records will do so. When bit 0 is true, all drivers generate or read unformatted records. For unformatted records the data in the external medium will be an exact image of the internal buffer. For formatted records certain drivers expect or supply additional formatting data. In most discussions bit 0 is referenced as "the special format bit" and unformatted records are called "special format records". Details of record formatting are described in the applicable driver documentation.

Bits 1 to 15 of word 6 are the End-Action Address (EAA). This address is the address of an end-action subroutine ^{link program} to which the I/O system will transfer when a particular I/O is completed. If this address is 0, no end action is specified. For a discussion of end action, refer to Section 1.6.

Word 7:

Word 7 is used by most drivers as temporary storage. It has a special use when the magnetic tape chain feature is present in the system, as described in the magnetic tape driver documentation.

1.4 BASIC IOS CALLS

IOS provides the user with 3 basic calls:

- a. OPEN
- b. DOIO
- c. STAT

With these three calls the user can perform an I/O operation on any standard 703 peripheral device.

In the discussion to follow calling sequence for the various I/O subroutine calls are stated in SYM II assembly language format.

```
CALL ARG1, ARG2...ARGN
```

This will be assembled as though it were written:

```
SMB      CALL
JSX      CALL
DATA     ARG1
DATA     ARG2
.....
DATA     ARGN + X'8000'
```

Note that the sign bit of the last argument is set. For more detail refer to the SYM II manual for an explanation of this form. Entry points to the I/O subroutines are through fixed link points in the System Linkage Table located in cells X'40' to X'7F'. These points must be defined to the assembler using equate cards. A listing of the System Linkage Table appears in Appendix D.

For example, the programmer who writes

```
DOIO  INFIOT
```

must include the definition

```
DOIO  EQU  X'44'
```

in his program. Programs written to execute on page 0 of the computer (locations 80_{16} through $7FF_{16}$) may omit the SMB instruction by using the form:

```
JSX   DOIO, INFIOT
```

Appendix A illustrates a sample input/output program which is written in the SYM I assembly language. The same program written in the SYM II assembly language is shown in Appendix B.

1.4.1 OPEN

The first call for any I/O operation is to subroutine OPEN. OPEN translates the arguments to the proper format and inserts them into the FIOT table.

OPEN has a fixed length calling sequence as follows:

```
OPEN  FIOT, BUF, WC, UNIT, OPER, MODE
```

where: FIOT is the first location of an 8-word array for the file description table.

BUF is the data buffer starting address.

WC is the address of the number of words.

UNIT is the logical unit number.

OPER is the type of operation (read, write, etc.).

MODE defines binary or alpha operation; 1 = binary; 0 = alpha.

It should be noted that the disk sector and the end-action address are not inserted into the FIOT by OPEN. These are less frequently used arguments and were eliminated to shorten the OPEN calling sequence. The user can assemble the arguments into the FIOT or store them at execution time.

The use of the OPEN call is optional, insofar as the FIOT may be assembled in the correct form. However, the OPEN call will initialize the FIOT, setting the busy bit off and clearing the status response word and may provide a convenient way of re-initializing. In addition, OPEN converts the specified logical unit into the actual physical unit number and stores it into bits 7 - 11 of word 2 of the FIOT, so that the using program may determine this assignment before beginning an I/O operation, if required.

1.4.2 DOIO

Actual I/O operations are initiated with calls to DOIO. DOIO changes the FIOT table as required, calls a driver to initiate the I/O operation, and returns to the user.

The calling sequence is as follows:

DOIO FIOT, BUF, WC

where: BUF is the data buffer starting address.

WC is the address of the number of words.

DOIO has a variable length calling sequence. If BUF or WC do not change from the previous call or have been set by OPEN, they need not be included. However, if WC is to be changed, BUF must be included.

1.4.3 STAT

STAT provides the user with a means of waiting until an I/O operation is complete. STAT has a variable length calling sequence as follows:

STAT FIOT, ERR

When a call to STAT is made, it will loop until the FIOT busy bit is off. If the ERR argument is not included, STAT will make a return to the next instruction following the calling sequence. If argument ERR is in the calling sequence, STAT will load the accumulator with the proper error processor return, and transfer to the ERR address. If the user wished to return from

to the ERR address. If the user wished to return from his ERR routine to the program sequence, he can do this by saving the accumulator as a return, then returning by loading the return into the index and executing a JMP*2 (see sample program in Appendices A and B). If there were no errors, STAT returns to the next instruction following the ERR argument.

1.5 SPECIAL IOS CALLS

The following calls perform special functions; such as Rewind, Write-End-of-File, and Backspace. Not all devices can perform the specified operation; for instance, the teletype will not rewind. If the device cannot perform the specified operation, the subroutine will return with the sign of the accumulator negative. The program can test this condition and determine an appropriate course of action. For example, the Loader accepts input from the BIN device (logical unit 4). If an input error occurs, the Loader backspaces the unit. If the unit was mag tape, it backspaces and the Loader re-reads the record. If the unit is not mag tape and cannot backspace, the Loader requests operator intervention to reposition the record.

The disk unit cannot perform the special functions; however, special calls for the disk return with the accumulator positive, since such functions as backspace and write end-of-file may be simulated by adjusting the sector address in the FIOT and adjusting the disk logical file pointers, which may be done arbitrarily and then allowing the program to drop through the special calls. Refer to the disk driver documentation.

The only information which must be in the FIOT for special operation is the logical unit number (bits 7 - 11 of word 2).

1.5.1 Device Status

The physical status of any I/O device may be interrogated at any time. The call is

STUS FIOT

The device's physical status response is returned in the accumulator. The FIOT is not disturbed by this operation, except to substitute the actual physical device number for the logical unit, if this has not already been done. An I/O operation in progress will not be disturbed.

1.5.2 Write End-of-File

A file mark is defined for most I/O devices and may be written on all output devices except the disk, line printer, and plotter. For magnetic tape the file mark is hardware supplied. For paper tape an ASCII "Bell" (X'87') character is the file mark. For cards a 2-7-8-9 punch is the file mark. All input devices upon sensing a file mark conclude the I/O operation and return End-of-File status (bit 10 of the status word true). The write end-of-file call is

WEOF FIOT

The FIOT will be busy until the operation is complete.

1.5.3 Seek End-of-File

A file mark is sought on the specified unit. Only the magnetic tape can perform this function. Refer to the magnetic tape driver documentation. The call is

SEOF FIOT

1.5.4 Backspace

The specified unit is backspaced one record. Only the magnetic tape can perform this operation. The call is

BKSP FIOT

1.5.5 Write-Skip

The specified magnetic tape unit will erase forward three inches of tape. The call is

WSKP FIOT

1.5.6 Rewind

The specified magnetic tape unit is rewound. The call is

REWD FIOT

1.6 END ACTION

The End Action feature permits program operation to be synchronized to I/O device operation. It allows sharing of time, otherwise wasted during I/O operations, with other tasks. A core resident program can initiate an I/O operation and then exit to allow the system to initiate another task. Then, when the I/O operation is complete, the system will return to the resident program at the end-action address. It can start another I/O operation and then exit. The system will restore the interrupted task.

End action may also be used within a single task; for example, to operate a double buffered I/O function.

End action is requested by specifying a non-zero service address in word 6 of the FIOT. Upon completion of the I/O operation, the I/O system will transfer to that address, after servicing but before restoring the final interrupt of the I/O operation. Then, when the end action subroutine exits, the system automatically restores the interrupted task.

End action subroutines are largely arbitrary in form and function. Any legal operation, including I/O, may be performed as part of an end-action subroutine. The system transfers to the end action address with the return on the index, so that the end action service routine may be in the form of a subroutine:

```

    STX    RETSAVE
    .
    .
    .
    LDX    RETSAVE
    JSX *  0
  
```

Of course, the SUBR and EXIT pseudo operations of SYM II may be used.

The end action service routine may also simply exit through the system EXIT link, cell X'40'. The system intercepts this exit and restores the interrupt. Thus, the service routine may have the form of a main program.

There are two restrictions on an end action service routine which affects the initiation of new I/O operations. When the new I/O operation is on the same device or on a device with the same interrupt level, as the device causing the end action, DOIO does not return after initiating the I/O, but rather exits the end action. This frees the interrupt level for the new I/O operation. Practically, this means that the last thing which should be done in an end action service routine is to initiate new I/O for the same device. It is emphasized that this restriction does not apply to I/O for devices on other interrupt levels, which return normally from the DOIO call.

The second restriction applies when the end action service routine is interrupting the interrupt service routine for the device for which the new I/O request is being made. When this occurs, the I/O in progress on that device cannot complete and, therefore, the new I/O operation cannot be initiated. In this event DOIO does not set the FIOT busy bit, but returns a status response of X'00FF' in word 3 of the FIOT. A STAT call with an error return will take the error return for this condition. This condition is virtually unrecoverable; however, it can occur normally only when two users are simultaneously attempting to use the same device. Cautious programming should avoid this situation. A possible solution to this situation, should it occur, would be to initiate a dummy I/O operation on an available device for example, the console typewriter whose end action would attempt to initiate the I/O on the unavailable device.

1.7 DEVICE DRIVER REFERENCES

<u>I. D. Code</u>	<u>Title</u>	<u>Drawing</u>
BPU	Disk and Magnetic Tape Driver	391040
BSU	Card Punch Driver	390018
BSV	Card Reader Driver	390019
BSW	Line Printer Driver	390020
BUU	Teletype Multiplexer Driver	391909
BYM	Teletype - High Speed Paper Tape I/O Driver	392292
BYP	Plotter Driver	392295

??
CK

APPENDIX A

SYM I SAMPLE PROGRAM

??
SA

* TEST PROGRAM TO READ A RECORD FROM
* THE ~~BINT~~^{DATA} UNIT AND TYPE OUT DONE.

ORIG 2048

READ EQU ~~95~~ * 'B'

WRIT EQU X'E *

PRIN BINT EQU 12

BIN EQU 1

ALPH EQU 0

LIST EQU 3

EXIT EQU 64

OPEN EQU 66

DOIO EQU 68

STAT EQU 70

0800 0080 STRT SMB OPEN

0801 2042 JSX OPEN,RFIT,RBUF,RWC,*PRIN* BINT,READ,*ALPH* BIN

* 0802 37FF

* 0803 37FF

* 0804 37FF

0805 0001

0806 0009

0807 8001

0808 0080 SMB OPEN

0809 2042 JSX OPEN,TFIT,RBUF,RWC,LIST,WRIT,ALPH

* 080A 37FF

* 080B 3003

* 080C 3004

080D 0003

080E 000E

080F 8000

0810 0080 SMB DOIO

0811 2044 JSX DOIO,RFIT

* 0812 3802

0813 0080 SMB STAT

0814 2046 JSX STAT,RFIT,RERR

* 0815 3012

* 0816 3FFF

APPENDIX A
(Continued)

SYM I SAMPLE PROGRAM

```

0817 0080      SMB  DOIO
0818 2044      JSX  DOIO,TFIT,FMES,TWO
* 0819 300A
* 081A 37FF
* 081B 3FFF

081C 0080      SMB  STAT
081D 2046      JSX  STAT,TFIT
* 081E 3819

081F 0080      SMB  EXIT
0820 2040      JSX  EXIT

          RBUF RES 10
          RFIT RES 8
0833 0002      TWO  D   2
* 0834 77FF      RERR STW RER
0835 0080      SMB  DOIO,TFIT
0836 2044      JSX  DOIO,TFIT
* 0837 381E
* 0838 9034      LDX  RER
0839 1802      JMP * 2
083A 0000      RER  DATA 0
          TFIT RES 8
0843 C4CF      FMES D   'DO','NE'
0844 CEC5
          VE
0845 000A      RWC  DATA 10

          END
    
```

APPENDIX A
(Continued)

SYM I SAMPLE PROGRAM

PSEUDO LISTING

0800	0080	STRT	0823	0000	
0801	2042		0824	0000	
0802	082B		0825	0000	
0803	0821		0826	0000	
0804	0846		0827	0000	
0805	0001		0828	0000	
0806	0009		0829	0000	
0807	8001		082A	0000	
0808	0080		082B	0000	RFIT
0809	2042		082C	0000	
080A	083B		082D	0000	
080B	0821		082E	0000	
080C	0846		082F	0000	
080D	0003		0830	0000	
080E	000E		0831	0000	
080F	8000		0832	0000	
0810	0080		0833	0002	TWO
0811	2044		0834	703A	RERR
0812	882B		0835	0080	
0813	0080		0836	2044	
0814	2046		0837	883B	
0815	082B		0838	903A	
0816	8834		0839	1802	
0817	0080		083A	0000	RER
0818	2044		083B	0000	IFIT
0819	083B		083C	0000	
081A	0843		083D	0000	
081B	8833		083E	0000	
081C	0080		083F	0000	
081D	2046		0840	0000	
081E	883B		0841	0000	
081F	0080		0842	0000	
0820	2040		0843	C4CF	FMES
0821	0000	RBUF	0844	CEC5	
0822	0000		0845	000A	RWC

APPENDIX B

SYM II SAMPLE PROGRAM

```

1 *      TEST PROGRAM TO READ A RECORD FROM
2 *      THE BINT UNIT AND TYPE OUT DONE.
3        ORIG  2048
0009     4 READ  EQU  9
000E     5 WRIT  EQU  X'E'
0001     6 BINT  EQU  1
0001     7 BIN   EQU  1
0000     8 ALPH  EQU  0
0003     9 LIST  EQU  3
0040    10 DONE  EQU  64
0042    11 OPEN  EQU  66
0044    12 DOIO  EQU  68
0046    13 STAT  EQU  70
S 0800 0080    14 STRT  OPEN  RFIT,RBUF,RWC,BINT,READ,BIN
S 0801 2042
0802 082B
0803 0821
0804 0844
0805 0001
0806 0009
0807 8001
0808 0080    15      OPEN  TFIT,RBUF,RWC,LIST,WRIT,ALPH
S 0809 2042
080A 083A
080B 0821
080C 0844
080D 0003
080E 000E
080F 8000
0810 0080    16      DOIO  RFIT
S 0811 2044
0812 882B
0813 0080    17      STAT  RFIT,RERR
S 0814 2046
0815 082B
0816 8833
0817 0080    18      DOIO  TFIT,FMES,=2
S 0818 2044
0819 083A
081A 0842
081B 8845
081C 0080    19      STAT  TFIT
S 081D 2046
081E 883A
081F 0080    20      DONE
S 0820 2040
0821        21 RBUF  RES  10
082B        22 RFIT  RES  8
0833 7039    23 RERR  STW  RER
0834 0080    24      DOIO  TFIT
S 0835 2044
0836 883A
0837 9039    25      LDX   RER
0838 1802    26      JMP  * 2
0839 0000    27 RER   D    0
083A        28 TFIT  RES  8

```


APPENDIX B
(Continued)
SYM II SAMPLE PROGRAM

```
0842 C4CF      29 FMES      TEXT  'DONE'  
0843 CEC5  
0844 000A      30 RWC       D      10  
          0800      31       END    STRT  
  
0845 0002
```

NO ERRORS

```
ALPH  0000  BIN      0001  BINT      0001  D010      0044  
DONE  0040  FMES     0842  LIST      0003  OPEN      0042  
RBUF  0821  READ     0009  RER       0839  RERR      0833  
RFIT  082B  RWC      0844  STAT      0046  STRT      0800  
TFIT  083A  WRIT     000E
```

NEXT

APPENDIX C

703 Character Codes

Internal Code	Graphic	Hollerith Card Code
A0	blank	blank
A1	!	11-2-8
A2	"	0-5-8
A3	#	0-7-8
A4	\$	11-3-8
A5	%	11-7-8
A6	&	12-2-8
A7	'	4-8
A8	(0-4-8
A9)	12-4-8
AA	*	11-4-8
AB	+	12
AC	,	0-3-8
AD	-	11
AE	.	12-3-8
AF	/	0-1
B0	0	0
B1	1	1
B2	2	2
B3	3	3
B4	4	4
B5	5	5
B6	6	6
B7	7	7
B8	8	8
B9	9	9
BA	:	5-8
BB	;	11-6-8
BC	<	12-6-8
BD	=	3-8
BE	>	6-8
BF	?	12-2-8

Internal Code	Graphic	Hollerith Card Code
C0	@	0-2-8
C1	A	12-1
C2	B	12-2
C3	C	12-3
C4	D	12-4
C5	E	12-5
C6	F	12-6
C7	G	12-7
C8	H	12-8
C9	I	12-9
CA	J	11-1
CB	K	11-2
CC	L	11-3
CD	M	11-4
CE	N	11-5
CF	O	11-6
D0	P	11-7
D1	Q	11-8
D2	R	11-9
D3	S	0-2
D4	T	0-3
D5	U	0-4
D6	V	0-5
D7	W	0-6
D8	X	0-7
D9	Y	0-8
DA	Z	0-9
DB	⌈	12-5-8
DC	⌋	0-6-8
DD	⌋	11-5-8
DE	↵	7-8
DF	↵	2-8

APPENDIX D
System Linkage Table

<u>Cell (Base 16)</u>	<u>Cell Use</u>
40	EXIT
42	OPEN
44	DOIO
46	STAT
48	BKSP
4A	WKSP
4C	SEOF
4E	RWND
50	WEOF
52	STUS
54	END Lower limit of available core (end X-RAY in Basic)
55	RBEG Restart address for RELO programs
56	Address of interrupt routine for TTY
57	ENB
58	STYP System type word
59	PEAT Address of system Peat Table
<u>5A</u>	ULIM Upper limit of available core*
5B	ENDM Upper or lower limit of Basic or Standard I/O Subsystems, respectively
<u>5C</u>	CORE Address of last cell in core
5D	DATE
61	NAME
65	JMP*0
66	LLIM Lowest cell loaded by loader
67	JOB Control Word
68	Unassigned
69	Unassigned
6A	QUEU Monitor Queue Subroutine
6C	QFCH Monitor Queue Fetch Argument Subroutine
6E	TYPE Type message subroutine
70	MTSF Masked test system flag word
72	LOAD Load a processor
74	DVEC Disk Vector Pointer
75	ULIM2 Upper limit of available core*
76	P. MAP Processor Map sector on disk
77	LOAD Load X-RAY Flag
78	ABSLOAD Absolute program load
7A	RLRS RELO program restart
7C	RESTL Resident task table linkage
7D	PFS Power Fail Safe Linkage
7F	RDTF Resident Disk Table size

*(Below resident task, blocks, etc.)

