

Enveloping an Iteration Scheme

Abstract: The problem of determining the optimal values of relaxation parameters for a linear iteration is solved. The optimal iteration scheme is achieved by a second-order linear iteration method.

Introduction

In this paper we discuss two procedures for accelerating iterative methods for solving the N th order linear system

$$Au = b, \quad (1)$$

where A is a given real nonsingular $N \times N$ matrix, and b (which is given) and the solution u are N -vectors. One procedure is the class of semi-iterative methods and the other is the class of relaxation methods. Both of these methods are described by R. S. Varga [1]. Both methods deal with an iteration scheme of the form

$$u_{n+1} = Gu_n + f. \quad (2)$$

The iteration scheme has the property that its fixed point v , where

$$v = Gv + f, \quad (3)$$

is the solution, u , of (1).

Both acceleration methods employ an additional set of parameters, one for each iteration for adjusting the value of the newly arrived at iterate. For a specified number of iterations, say m , we consider the problem of optimally determining these parameters, i. e., determining them so that the error between the m th iterate and the exact solution is minimized. For semi-iterative methods this problem has been solved [2] and the parameters are characterized as the coefficients of an appropriately normalized Chebyshev polynomial of degree m . In this paper we show that the optimization problem for the class of relaxation methods is solved by choosing parameters which are the roots of an appropriately determined m th degree orthogonal polynomial.

For semi-iterative methods, the optimal iteration scheme is achievable by an appropriate iteration method of second degree, i.e., one of the form

$$v_{n+1} = (\alpha_n G + \beta_n)v_n + \gamma_n v_{n-1} + k_n. \quad (4)$$

By this we mean that for appropriate choices of α_n , β_n , γ_n , and k_n , v_n coincides with the n th iterate of the optimal semi-iterative method of n steps, for any n .

We will show that the same is true for the optimal relaxation method. We refer to this method of second degree as the *envelope* of the class of relaxation methods. We use this terminology for the following reason. For a given choice, say $u_0 = 0$, consider the polygonal path joining the points $(i, \|u_i - u\|)$, $i = 0, 1, \dots, n$, where $\|\cdot\|$ denotes the norm of the vector u . This polygonal path characterizes the iteration scheme for each n and each choice of n relaxation parameters. The totality of all such polygonal paths corresponding to all possible relaxation schemes has a lower envelope that is yet another polygonal path. This envelope is described by the second-degree method.

Our solution consists in adapting to this more general case a special case that has been solved by E. Stiefel [3], where $Gu + k$ is the gradient of a quadratic form. In this case the iteration scheme is a relaxation version of the method of steepest descent and the envelope is the method of conjugate gradients.

The optimal method

Consider the linear system

$$Au = b. \quad (1)$$

Let the iteration scheme

$$u_{n+1} = Gu_n + f, \quad n = 0, 1, \dots \quad (5)$$

where G is an $N \times N$ matrix and f an N -vector have a unique fixed point that coincides with the solution of (1). We write (5) as

$$\Delta u_\nu = r_\nu, \quad \nu = 0, 1, \dots, \quad (6)$$

where

$$\Delta u_\nu = u_{\nu+1} - u_\nu, \quad (7)$$

$$r_\nu = f - Mu_\nu, \text{ and}$$

$$M = I - G.$$

where I is the identity matrix.

We now introduce the first acceleration procedure for (5).

• Semi-iterative methods

Let $\alpha_{n,k}$, $k = 0, 1, \dots, n$, $n = 0, 1, \dots$ be a sequence of real constants with the property

$$\sum_{k=0}^n \alpha_{n,k} = 1, \quad n = 0, 1, \dots \quad (8)$$

and let

$$v_n = \sum_{k=0}^n \alpha_{n,k} u_k, \quad n = 0, 1, \dots \quad (9)$$

The v_n define a semi-iterative method of successive approximations to u . The error $e_n \equiv v_n - u$ satisfies

$$e_n = P_n(G)e_0, \quad n = 1, 2, \dots, \quad (10)$$

$$P_n(x) = \sum_{k=0}^n \alpha_{n,k} x^k. \quad (11)$$

in which $P_n(1) \equiv 1$.

Let the eigenvalues λ of G be real and suppose that they lie in the interval $[\alpha, \beta]$ with $\alpha < \beta < 1$. Then for the eigenvalues $P_n(\lambda)$ of $P_n(G)$, we see that the

$$\max_{\alpha \leq \lambda \leq \beta} |P_n(\lambda)| \quad (12)$$

is minimized under the constraint $P_n(1) = 1$ for

$$P_n(x) = T_n\left(\frac{2x - (\beta + \alpha)}{\beta - \alpha}\right) / T_n(\gamma), \quad (13)$$

where

$$\gamma = \frac{2 - \beta - \alpha}{\beta - \alpha} \quad (14)$$

and the T_n are the Chebyshev polynomials of degree n defined by

$$T_n(x) = \cos(n \cos^{-1} x), \quad n = 0, 1, \dots, \quad (15)$$

if $|x| \leq 1$.

Moreover from the property

$$T_{n+1}(x) = 2xT_n(x) - T_{n-1}(x) \quad (16)$$

of the Chebyshev polynomials, it can be deduced that the v_n of Eq. (9) for the choice of P_n in (13) obey the following second-order recurrence relation:

$$v_{n+1} = (\alpha_n G + \beta_n I)v_n + \gamma_n v_{n-1} + f_n, \quad n = 1, 2, \dots, \quad (17)$$

where

$$\alpha_n = \frac{2}{(\beta - \alpha)\gamma w_{n+1}}$$

$$\beta_n = \frac{\beta + \alpha}{(\alpha - \beta)\gamma w_{n+1}}$$

$$\gamma_n = 1 - w_{n+1}$$

$$f_n = \frac{2 w_{n+1}}{\gamma (\beta - \alpha)} f$$

$$w_1 = 1$$

$$w_2 = \frac{1}{1 - 1/2\gamma^2}$$

$$w_{n+1} = \frac{1}{1 - w_n/4\gamma^2}, \quad n = 2, 3, \dots \quad (18)$$

The choice of the $\alpha_{n,k}$ gives the optimal acceleration to the semi-iterative scheme and the property (16) enables one to avoid using the coefficients of the Chebyshev polynomials and resumming for every change in n , by use of the second-order method (17).

Now let us consider the second acceleration procedure.

• Relaxation methods

Let t_n , $n = 0, 1, \dots$ be a sequence of positive constants. Let x_n , $n = 0, 1, \dots$ define a relaxed iterative method corresponding to (5), where

$$x_{n+1} - x_n = (1/t_n)(f - Mx_n). \quad (19)$$

We write this as

$$\Delta x_n = r_n/t_n, \quad (20)$$

where $\Delta x_n = x_{n+1} - x_n$ and $r_n = f - Mx_n$.

The error $e_n = x_n - u$ satisfies

$$e_n = -M^{-1}r_n, \quad (21)$$

if M is nonsingular.

Suppose that M is positive definite and Hermitian. Then if $(u, v) = u^*v$ denotes the usual inner product in N -space, (u, Mu) defines a metric in this space which we denote by $\mu(u)$. In particular for the N -vector $e = -M^{-1}r$, we have

$$\mu(e) = (M^{-1}r, r) \quad (22)$$

For each n , we consider the problem of choosing the

$t_j, j = 0, 1, \dots, n$ so that $\mu(e_n)$ is a minimum. Of course the optimal relaxation parameters will depend on n , but for convenience we suppress displaying this dependence. Following Stiefel, we may formulate the solution to this problem as follows.

Let $R_n(\lambda), n = 1, 2, \dots$, be a sequence of polynomials given by

$$R_n(\lambda) = \prod_{j=1}^n (1 - \lambda/t_{j-1}). \quad (23)$$

If the initial iterate is $x_0 = 0$, then

$$x_1 = \frac{f}{t_0} = M^{-1} \left[I - \left(I - \frac{M}{t_0} \right) \right] f = M^{-1} [I - R_1(M)] f.$$

Then

$$x_n = M^{-1} [I - R_n] f, \quad (24)$$

since by induction

$$\begin{aligned} x_{n+1} &= M^{-1} \left[I - R_n + \frac{M}{t_n} - \frac{M}{t_n} (I - R_n) \right] f \\ &= M^{-1} [I - R_{n+1}] f. \end{aligned}$$

Similarly

$$r_n = R_n(M) f \quad (25)$$

so that from (22)

$$\mu(e_n) = (M^{-1} R_n(M) f, R_n(M) f). \quad (26)$$

Let M have distinct eigenvalues $0 < \lambda_1 < \lambda_2 < \dots < \lambda_N < 1$ and let T be the matrix whose j th column is the eigenvector of M with eigenvalue $\lambda_j, j = 1, \dots, N$. Then

$$\begin{aligned} \mu(e_n) &= (M^{-1} T T^{-1} R_n(M) T T^{-1} f, T T^{-1} R_n(M) T T^{-1} f) \\ &= (\Lambda^{-1} R_n(\Lambda) \bar{f}, R_n(\Lambda) \bar{f}). \end{aligned} \quad (27)$$

Here Λ is the $N \times N$ diagonal matrix, $\Lambda = (\lambda_j \delta_{ij})$ and $\bar{f} = T^{-1} f$.

Then

$$\mu(e_n) = \sum_{j=1}^N \bar{f}_j^2 [R_n^2(\lambda_j)] / \lambda_j, \quad (28)$$

where $\bar{f}_j, j = 1, \dots, N$ are the components of $\bar{f}$. Thus if we introduce the density

$$\rho(\lambda) = \sum_{j=1}^N \bar{f}_j^2 \delta(\lambda - \lambda_j), \quad (29)$$

we may write

$$\mu(e_n) = \int_0^1 [\rho(\lambda)/\lambda] R_n^2(\lambda) d\lambda. \quad (30)$$

In the case that the eigenvalues of M are not distinct this representation for $\mu(e_n)$ will in general involve $R_n(\lambda)$ and some of its derivatives. We do not treat this more involved case further.

Now the problem of optimal choice of the relaxation parameters corresponds to finding the polynomial $R_n(\lambda)$ which solves the polynomial least-squares problem with respect to the weight $\rho(\lambda)/\lambda$. Of course we must impose the additional constraint

$$R_n(0) = 1, \quad n = 0, 1, \dots$$

if we are to recover polynomials of the type (23). Denoting the class of polynomial of degree n as Π_n , the problem to be solved is explicitly

$$\min_{\substack{R_n \in \Pi_n \\ R(0)=1}} \mu(e_n) = \min_{\substack{R_n \in \Pi_n \\ R(0)=1}} \int_0^1 [\rho(\lambda)/\lambda] R_n^2(\lambda) d\lambda.$$

The solution to this constrained least-squares problem is the set of polynomials $P_n(\lambda)$, orthonormal on $(0,1)$ with respect to the weight $\rho(\lambda)$ (see [4]), i.e.,

$$\int_0^1 P_\mu(\lambda) P_\nu(\lambda) \rho(\lambda) d\lambda = 0, \quad \mu \neq \nu. \quad (31)$$

Thus the optimal relaxation parameters are $t_{j,n}, j = 1, \dots, n, n = 1, 2, \dots$, where for each n , the $t_{j,n}$ are the roots of $P_n(\lambda)$.

As in the case for the optimal semi-iterative method, the optimal relaxation method may be replaced by a second-order iteration method that frees the method from use of the roots $t_{j,n}$ as well as recomputing for each n change. This second-order method emerges from the second-order recurrence relation

$$\begin{aligned} \lambda R_i(\lambda) &= -q_i R_{i+1}(\lambda) + (p_i + q_i) R_i(\lambda) - p_i R_{i-1}(\lambda), \\ i &= 1, 2, \dots \end{aligned} \quad (32)$$

obeyed by the $R_i(\lambda)$, which we henceforth assume to satisfy (31). We take

$$R_0 = 1, p_0 = 0, q_{-1} = 0,$$

and

$$\begin{aligned} q_i &= (1/N_i) (r_i, M r_i) - p_i \\ &= (1/N_i) \int_0^1 \lambda R_i^2(\lambda) \rho(\lambda) d\lambda - p_i \end{aligned}$$

$$p_i = (N_i/N_{i-1}) q_{i-1}$$

$$N_i = (r_i, r_i) = \int_0^1 R_i^2(\lambda) \rho(\lambda) d\lambda, \quad i = 0, 1, 2, \dots \quad (33)$$

Note that the p_i and q_i may be computed without any use of the orthogonal polynomials.

The second-order scheme is

$$\Delta z_i = (1/q_i) (r_i + p_i \Delta z_{i-1}), \quad z_0 = 0. \quad (34)$$

We claim that this scheme envelopes the optimal relaxation scheme for each n , i.e., $z_i = M^{-1} [I - R_i(M)] f$ [see (24)]. This may be seen as follows.

Since $z_0 = 0$, $r_0 = f = R_0(M)f$. Suppose by induction that $r_i = R_i(M)f$. Now since $r_i = f - Mz_i$, we have

$$r_{i+1} - r_i = -M\Delta z_i. \quad (35)$$

Then, from Eq. (34),

$$\begin{aligned} r_{i+1} - r_i &= -\frac{M}{q_i}(r_i + p_i \Delta z_{i-1}) \\ &= -\frac{Mr_i}{q_i} + \frac{p_i}{q_i}(r_i - r_{i-1}) \\ &= \left[\left(\frac{p_i}{q_i} - \frac{M}{q_i} \right) R_i(M) - \frac{p_i}{q_i} R_{i-1}(M) \right] f. \end{aligned} \quad (36)$$

Using (32) this becomes

$$r_{i+1} = R_{i+1}(M)f. \quad (37)$$

Then having completed the induction we may write

$$z_i = M^{-1}(f - r_i) = M^{-1}[I - R_i(M)]f, \quad (38)$$

which is what was to be proved.

Remark: Notice that the three-step recurrence scheme (34) must theoretically terminate after N steps. This follows from the fact that since $\rho(\lambda)$ as given in Eq. (29) has N jumps, then only N orthogonal polynomials $P_n(\lambda)$ can be generated by (31).*

Examples

From $Au = b$, certain iteration schemes are obtained by splitting A

$$A = P - N \quad (39)$$

where N is nonsingular. Then writing

$$Nu_{n+1} = Pu_n - b. \quad (40)$$

*The author is grateful to R. Varga for urging him to include this remark.

Then

$$\begin{aligned} f &= -N^{-1}b \\ G &= N^{-1}P \\ M &= I - N^{-1}P \end{aligned} \quad (41)$$

- 1) The Jacobi method corresponds to $N = \text{diag } A$.
- 2) The Gauss-Seidel method corresponds to $N = L$ (which equals the lower triangular portion of A up to and including the main diagonal, with all other elements of L taken as zero).
- 3) The method of steepest descent corresponds to $M = -A$ and $f = b$.

In each of these three cases we must take care that the associated matrix M has the properties required in the section on the optimal method, i.e., Hermitian and positive definite.

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